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Cooperative Data Protection in Topology-Aware Decentralized Storage Networks

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Outline

Introduction

- Motivation and Model
- Existing work

Cooperative Data Protection

- ECC hierarchy
- Single-level cooperation

Multi-level cooperation

- Cooperation graphs and compatible graphs
- Construction over compatible graphs

Conclusion

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Motivation

- Invention of Blockchain technology makes the concept of "decentralization" popular
 - Higher privacy
 - Better scalability and flexibility
 - Economically attractive



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- Invention of Blockchain technology makes the concept of "decentralization" popular
 - Higher privacy
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 - Economically attractive
- Decentralization has potential to universally revolutionize various applications
 - Decentralized storage networks
 - Masterless coded distributed computation
 - Federated learning
 - Wireless sensor networks



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Dynamic DSNs: Structures

- DSN: distributed local clouds
 - Clusters of nodes
 - Local messages are encoded and stored distributively among local nodes



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Dynamic DSNs: Practical Concerns



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Desired Properties: Reliability and Latency

Hierarchical erasure correction

- Each node provides different levels of robustness for the codeword stored at it through accessing different sets of nodes
 - Trade-off between reliability and latency

Topology-awareness

- Schemes optimized for DSNs with a specific topology can result in bad performance in DSNs with other topologies
 - Intrinsic heterogeneity
 - Latency for inter-cloud communication can be much higher than that of intra-cloud communication



Desired Properties: Scalability and Flexibility

Scalability

- To support node churn in DSN
- Expand the backbone network to accommodate additional workload without rebuilding the entire infrastructure

Flexibility

- To support the dynamic nature of the usage rate of data
- Split a local cloud into smaller clouds if the data stored in it become hot

Our goal is to construct topology-aware coding schemes that provide hierarchical erasure correction at each node and simultaneously support scalability and flexibility

Abstract Model: Parameters

- A DSN is modeled as a graph G(V, E)
 - Number of master nodes: p
 - > Message (codeword) stored at v_i : $\mathbf{m}_i(\mathbf{c}_i)$; lengths $k_i(n_i)$
 - > Redundancy of \mathbf{c}_i : $r_i = n_i k_i$
 - $\succ \mathbf{m} = (\mathbf{m}_1, \mathbf{m}_2, \cdots, \mathbf{m}_p)$
 - $\succ \mathbf{c} = (\mathbf{c}_1, \mathbf{c}_2, \cdots, \mathbf{c}_p)$



Existing Literature

Distributed storage^[1]

> No explicit consideration of clustering nature of network nodes

Multi-rack storage^[2-8]

- Network topologies are predetermined
- Capacities of the communication links are typically considered to be the same

^[1] A. G. Dimakis et al., "Network coding for distributed storage systems", IEEE Trans. Inf. Theory, vol. 56, no. 9, pp. 4539-4551, 2010

^[2] Z. Kong et al., "Decentralized coding algorithms for distributed storage in wireless sensor networks", IEEE JSAC, vol. 28, no. 2, pp. 261-267, 2010

^[3] M. Ye et al., "Cooperative repair: Constructions of optimal MDS codes for all admissible parameters", IEEE Trans. Inf. Theory, vol. 65, no. 3, pp. 1639-1656, 2018

^[4] N. Prakash et al., "The storage versus repair-bandwidth trade-off for clustered storage systems", IEEE Trans. Inf. Theory, vol. 64, no. 8, pp. 5783-5805, 2018

^[5] J. Li et al., "Tree-structured data regeneration in distributed storage systems with regenerating codes", IEEE INFOCOM, 2010

^[6] Y. Wang et al., "Non-homogeneous two-rack model for distributed storage systems", IEEE INFOCOM, 2014

^[7] H. Hou et al., "Rack-aware regenerating codes for data centers", IEEE Trans. Inf. Theory, 2019

^[8] Z. Chen et al., "Explicit constructions of MSR codes for clustered distributed storage: the rack-aware storage model", [Online]. Available: https://arxiv.org/abs/1901.04419, 2019

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ECC hierarchy describes the erasure correction (EC) capabilities of nodes while cooperating with different sets of other nodes



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- Cooperation at each node
 - > ECC hierarchy at v_i : $\mathbf{d}_i = (d_{i,0}, d_{i,1}, \dots, d_{i,L_i})$

•
$$\emptyset \subset \mathcal{A}_i^1 \subset \mathcal{A}_i^2 \subset \cdots \subset \mathcal{A}_i^{L_i} \subseteq V, \left\{\mathcal{B}_i^l\right\}_{1 \leq l \leq L_i}, \mathcal{A}_i^1 \cap \mathcal{B}_i^l = \emptyset$$

•
$$d_{i,l}$$
 : EC capability at v_i if nodes in $\mathcal{A}_i^l \cup \mathcal{B}_i^l$ are recovered



ECC hierarchy describes the erasure correction (EC) capabilities of nodes while cooperating with different sets of other nodes

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- $d_{i,l}$: EC capability at v_i if nodes in \mathcal{A}_i^l are recovered
- > Finer EC capability: $d_{i,l} = (\lambda_{i,l;\mathcal{W}})_{\emptyset \subseteq \mathcal{W} \subseteq \mathcal{B}_i^l}$
 - λ_{i,l;W}: EC capability while nodes in W are also recovered and involved in the *l*-th level cooperation at v_i

 $d_{i,0}$

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 $d_{i,3}$

Example

> Local decoding: $d_{i,0}$



Example

- > Local decoding: $d_{i,0}$
- $\succ v_2$ obtains extra parities from v_1

•
$$\lambda_{2,1;\mathcal{W}} = \lambda_{2,1;\emptyset}, \mathcal{W} \subseteq \{v_6, v_8\}$$



Example

- > Local decoding: $d_{i,0}$
- $\succ v_2$ obtains extra parities from v_1
 - $\lambda_{2,1;\mathcal{W}} = \lambda_{2,1;\emptyset}, \mathcal{W} \subseteq \{v_6, v_8\}$
- $\succ v_2$ obtains extra parities from v_3 if v_4 is also recovered
 - $\lambda_{2,1;\mathcal{W}} = \lambda_{2,1;\{v_4\}}, \{v_4\} \subseteq \mathcal{W} \subset \{v_4, v_6, v_8\}$



Example

- \succ Local decoding: $d_{i,0}$
- $\succ v_2$ obtains extra parities from v_1
 - $\lambda_{2,1:\mathcal{W}} = \lambda_{2,1:\emptyset}, \mathcal{W} \subseteq \{v_6, v_8\}$
- \succ v_2 obtains extra parities from v_3 if v_4 is also recovered
 - $\lambda_{2,1;W} = \lambda_{2,1;\{v_4\}}, \{v_4\} \subseteq W \subset \{v_4, v_6, v_8\}$
- $\succ v_2$ obtains extra parities from v_5 if $\{v_4, v_6, v_8\} \subseteq \mathcal{W}$

•
$$\lambda_{2,1;\mathcal{W}} = \lambda_{2,1;\{v_4,v_6,v_8\}} = d_{2,1}, \mathcal{W} = \{v_4, v_6, v_8\}$$



 Parity part of the generator matrix of a single-level accessible code based on CRS codes

$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	0	0	0	0	0	0	0	0	0	0
$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_9$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_8$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	${f A}_{10,10}$	$\mathbf{B}_{10,11}\mathbf{U}_{11}$	$\mathbf{B}_{10,12}\mathbf{U}_{12}$
0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	${f A}_{11,11}$	$\mathbf{B}_{11,12}\mathbf{U}_{12}$
0	0	0	0	0	0	0	0	0	$\mathbf{B}_{12,10}\mathbf{U}_{10}$	${f B}_{12,11}{f U}_{11}$	${f A}_{12,12}$



 Parity part of the generator matrix of a single-level accessible code based on CRS codes

-												
	$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	0	0	0	0	0	0	0	0	0	0
	$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
	0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
	0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
	0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
	0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
	0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_{9}$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
	0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
	0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_8$	$\mathbf{A}_{9,9}$	${f B}_{9,10}{f U}_{10}$	0	0
	0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	$\mathbf{A}_{10,10}$	${f B}_{10,11} {f U}_{11}$	$\mathbf{B}_{10,12}\mathbf{U}_{12}$
	0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	$\mathbf{A}_{11,11}$	${f B}_{11,12}{f U}_{12}$
	0	0	0	0	0	0	0	0	0	$\mathbf{B}_{12,10}\mathbf{U}_{10}$	${f B}_{12,11}{f U}_{11}$	$\mathbf{A}_{12,12}$
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Diagonal components: A_{i,i}



 Parity part of the generator matrix of a single-level accessible code based on CRS codes

$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	0	0	0	0	0	0	0	0	0	0
$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2.5}\mathbf{U}_5$	0	0	0	0	0	0	0
0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_{8}$	0	0	0	0
0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_{9}$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_{9}$	0	0	0
0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_{8}$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	${f A}_{10.10}$	$\mathbf{B}_{10,11}\mathbf{U}_{11}$	${f B}_{10,12} {f U}_{12}$
0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	${f A}_{11,11}$	${f B}_{11,12}{f U}_{12}$
0	0	0	0	0	0	0	0	0	$\mathbf{B}_{12,10}\mathbf{U}_{10}$	$\mathbf{B}_{12,11}\mathbf{U}_{11}$	$\mathbf{A}_{12,12}$

- > Diagonal components: $A_{i,i}$
- Non-diagonal components: B_{i,j}U_j



 Parity part of the generator matrix of a single-level accessible code based on CRS codes

$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	0	0	0	0	0	0	0	0	0	0
$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_9$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_8$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	${f A}_{10,10}$	${f B}_{10,11} {f U}_{11}$	${f B}_{10,12} {f U}_{12}$
0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	$A_{11,11}$	${f B}_{11,12}{f U}_{12}$
0	0	0	0	0	0	0	0	0	$\mathbf{B}_{12,10}\mathbf{U}_{10}$	${f B}_{12,11}{f U}_{11}$	$\mathbf{A}_{12,12}$

- > Diagonal components: $\mathbf{A}_{i,i}$ ($k_i \times r_i$)
- > Non-diagonal components: $\mathbf{B}_{i,j}\mathbf{U}_j(k_i \times \delta_j, \delta_j \times r_j)$
- > Components are parts of Cauchy matrices T_i

$$\mathbf{T}_i = egin{bmatrix} \mathbf{A}_{i,i} & \mathbf{B}_{i,j_1} & \dots & \mathbf{B}_{i,j_{|\mathcal{M}_i|}} \ \hline \mathbf{U}_i & \mathbf{Z}_i \end{bmatrix}$$

• \mathcal{M}_i : nodes cooperating with v_i in the 1-st level cooperation

$$d_{i,0} = r_i - \delta_i, \, d_{i,1} = r_i + \sum_{\nu_j \in \mathcal{M}_i} \delta_j, \, \lambda_{i,1;\mathcal{W}} = r_i + \sum_{\mathcal{M}_j \setminus \{\nu_i\} \subseteq \mathcal{M}_i \cup \mathcal{W}} \delta_j$$

Cauchy Matrices: Scalability and Flexibility

Component matrices

$$\mathbf{T}_i = egin{bmatrix} \mathbf{A}_{i,i} & \mathbf{B}_{i,j_1} & \dots & \mathbf{B}_{i,j_{|\mathcal{M}_i|}} \ \hline \mathbf{U}_i & \mathbf{Z}_i \end{bmatrix}$$

$$d_{i,0} = r_i - \delta_i, \, d_{i,1} = r_i + \sum_{v_j \in \mathcal{M}_i} \delta_j, \, \lambda_{i,1;\mathcal{W}} = r_i + \sum_{\mathcal{M}_j \setminus \{v_i\} \subseteq \mathcal{M}_i \cup \mathcal{W}} \delta_j$$

• Why Cauchy matrices?

- Component matrices at different nodes can be chosen independently – facilitates scalability
- Cauchy-like structure is inherited through concatenation and splitting of Cauchy matrices
 - Concatenating with a Cauchy matrix of dimension k corresponds to adding k extra parities
 - Splitting the component matrices corresponds to splitting a local cloud
 - -- facilitates flexibility

Example: Hierarchical Erasure Correction

	DIT	0	0	0	0	0	0	0	0	0	0
$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	U	U	U	U	U	U	U	U	U	U
$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_{9}$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_8$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	$A_{10,10}$	$\mathbf{B}_{10,11}\mathbf{U}_{11}$	$\mathbf{B}_{10,12}\mathbf{U}_{12}$
0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	$\mathbf{A}_{11,11}$	$\mathbf{B}_{11,12}\mathbf{U}_{12}$
0	0	0	0	0	0	0	0	0	$\mathbf{B}_{12,10}\mathbf{U}_{10}$	${f B}_{12,11}{f U}_{11}$	${f A}_{12,12}$

- ECC hierarchy at node v₂
 - > Local decoding: $d_{2,0} = r_2 \delta_2$



Example: Hierarchical Erasure Correction

$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	0	0	0	0	0	0	0	0	0	0
$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_9$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_8$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	$A_{10,10}$	$\mathbf{B}_{10,11}\mathbf{U}_{11}$	$\mathbf{B}_{10,12}\mathbf{U}_{12}$
0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	${f A}_{11,11}$	$\mathbf{B}_{11,12}\mathbf{U}_{12}$
0	0	0	0	0	0	0	0	0	$\mathbf{B}_{12,10}\mathbf{U}_{10}$	${f B}_{12,11}{f U}_{11}$	${f A}_{12,12}$

• ECC hierarchy at node v_2

- > Local decoding: $d_{2,0} = r_2 \delta_2$
- Cross parities at node v_2 can be subtracted from its local parity part if v_1, v_3, v_5 are recovered



Example: Hierarchical Erasure Correction

$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	0	0	0	0	0	0	0	0	0	0
$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_9$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_8$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	$A_{10,10}$	$\mathbf{B}_{10,11}\mathbf{U}_{11}$	$\mathbf{B}_{10,12}\mathbf{U}_{12}$
0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	${f A}_{11,11}$	${f B}_{11,12}{f U}_{12}$
0	0	0	0	0	0	0	0	0	$\mathbf{B}_{12,10}\mathbf{U}_{10}$	${f B}_{12,11}{f U}_{11}$	${f A}_{12,12}$

ECC hierarchy at node v₂

- > Local decoding: $d_{2,0} = r_2 \delta_2$
- Cross parities at node v_2 can be subtracted from its local parity part if v_1, v_3, v_5 are recovered
- Additional cross parities of c₂ can be obtained from v_j if all other neighbors of v_j except for v₂ are recovered:

•
$$\lambda_{2,1;\mathcal{W}} = r_2 + \sum_{\mathcal{M}_j \setminus \{v_2\} \subseteq \mathcal{M}_i \cup \mathcal{W}} \delta_j$$



$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	0	0	0	0	0	0	0	0	0	0
$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_9$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_8$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	$A_{10,10}$	${f B}_{10,11} {f U}_{11}$	$\mathbf{B}_{10,12}\mathbf{U}_{12}$
0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	$\mathbf{A}_{11,11}$	$\mathbf{B}_{11,12}\mathbf{U}_{12}$
0	0	0	0	0	0	0	0	0	$\mathbf{B}_{12,10}\mathbf{U}_{10}$	${f B}_{12,11}{f U}_{11}$	$\mathbf{A}_{12,12}$

- Suppose c_2 has $(r_2 + \delta_1 + 1)$ erasures
 - \succ Needs to obtain additional cross parities from v_3 or v_5
- Suppose transmission over e_{i,j} needs time t_{i,j}



$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	0	0	0	0	0	0	0	0	0	0
$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_9$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_8$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	$A_{10,10}$	${f B}_{10,11} {f U}_{11}$	$\mathbf{B}_{10,12}\mathbf{U}_{12}$
0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	$\mathbf{A}_{11,11}$	$\mathbf{B}_{11,12}\mathbf{U}_{12}$
0	0	0	0	0	0	0	0	0	$\mathbf{B}_{12,10}\mathbf{U}_{10}$	${f B}_{12,11}{f U}_{11}$	$\mathbf{A}_{12,12}$

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 - \succ Time cost through v_3

•
$$T_3 = t_{4,3} + t_{3,2}$$



$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	0	0	0	0	0	0	0	0	0	0
$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_9$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_8$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	$\mathbf{A}_{10,10}$	$\mathbf{B}_{10,11}\mathbf{U}_{11}$	$\mathbf{B}_{10,12}\mathbf{U}_{12}$
0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	$A_{11,11}$	$\mathbf{B}_{11,12}\mathbf{U}_{12}$
0	0	0	0	0	0	0	0	0	${f B}_{12,10}{f U}_{10}$	${f B}_{12,11}{f U}_{11}$	$\mathbf{A}_{12,12}$

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$$T_5 = \max_{i=4,6,8} t_{i,5} + t_{5,2}$$



$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	0	0	0	0	0	0	0	0	0	0
$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_9$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_{8}$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	$A_{10,10}$	$\mathbf{B}_{10,11}\mathbf{U}_{11}$	$\mathbf{B}_{10,12}\mathbf{U}_{12}$
0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	${f A}_{11,11}$	$\mathbf{B}_{11,12}\mathbf{U}_{12}$
0	0	0	0	0	0	0	0	0	${f B}_{12,10}{f U}_{10}$	$\mathbf{B}_{12,11}\mathbf{U}_{11}$	$\mathbf{A}_{12,12}$

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A	BUIL	0	0	0	0	0	0	0	0	0	0
$A_{1,1}$	$\mathbf{D}_{1,2}\mathbf{O}_2$	0	- 0	0		0	0	0	0	0	0
$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_8$	$\mathbf{B}_{7,9}\mathbf{U}_9$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_8$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	$A_{10,10}$	$\mathbf{B}_{10,11}\mathbf{U}_{11}$	$\mathbf{B}_{10,12}\mathbf{U}_{12}$
0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	${f A}_{11,11}$	$\mathbf{B}_{11,12}\mathbf{U}_{12}$
0	0	0	0	0	0	0	0	0	${f B}_{12,10}{f U}_{10}$	${f B}_{12,11}{f U}_{11}$	$\mathbf{A}_{12,12}$

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Nodes in the DSN automatically choose the fastest path for recovery



- Error pattern: v_2 , v_4 , v_8 , v_{10} are not locally recoverable
 - Recoverable in topology-aware coding
 - Each red node cooperates with 3 nodes



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Topology-aware coding enables information flow from more reliable nodes to less reliable nodes, tolerating more flexible erasure patterns



Outline

- Introduction
 - Motivation and Model
 - Existing work
- Cooperative Data Protection
 - ECC hierarchy
 - Single-level cooperation
- Multi-level cooperation
 - Cooperation graphs and compatible graphs
 - Construction over compatible graphs
- Conclusion

Parity part of the generator matrix

_												
	$\mathbf{A}_{1,1}$	$\mathbf{B}_{1,2}\mathbf{U}_2$	0	0	0	0	0	0	0	0	0	0
	$\mathbf{B}_{2,1}\mathbf{U}_1$	$\mathbf{A}_{2,2}$	$\mathbf{B}_{2,3}\mathbf{U}_3$	0	$\mathbf{B}_{2,5}\mathbf{U}_5$	0	0	0	0	0	0	0
	0	$\mathbf{B}_{3,2}\mathbf{U}_2$	$\mathbf{A}_{3,3}$	$\mathbf{B}_{3,4}\mathbf{U}_4$	0	0	0	0	0	0	0	0
	0	0	$\mathbf{B}_{4,3}\mathbf{U}_3$	$\mathbf{A}_{4,4}$	$\mathbf{B}_{4,5}\mathbf{U}_5$	$\mathbf{B}_{4,6}\mathbf{U}_{6}$	0	0	0	0	0	0
	0	$\mathbf{B}_{5,2}\mathbf{U}_2$	0	$\mathbf{B}_{5,4}\mathbf{U}_4$	$\mathbf{A}_{5,5}$	$\mathbf{B}_{5,6}\mathbf{U}_{6}$	0	$\mathbf{B}_{5,8}\mathbf{U}_8$	0	0	0	0
	0	0	0	$\mathbf{B}_{6,4}\mathbf{U}_4$	$\mathbf{B}_{6,5}\mathbf{U}_5$	$\mathbf{A}_{6,6}$	$\mathbf{B}_{6,7}\mathbf{U}_7$	0	0	0	0	0
	0	0	0	0	0	$\mathbf{B}_{7,6}\mathbf{U}_{6}$	$\mathbf{A}_{7,7}$	$\mathbf{B}_{7,8}\mathbf{U}_{8}$	$\mathbf{B}_{7,9}\mathbf{U}_{9}$	0	$\mathbf{B}_{7,11}\mathbf{U}_{11}$	0
_	0	0	0	0	$\mathbf{B}_{8,5}\mathbf{U}_5$	0	$\mathbf{B}_{8,7}\mathbf{U}_7$	$\mathbf{A}_{8,8}$	$\mathbf{B}_{8,9}\mathbf{U}_9$	0	0	0
	0	0	0	0	0	0	$\mathbf{B}_{9,7}\mathbf{U}_7$	$\mathbf{B}_{9,8}\mathbf{U}_8$	$\mathbf{A}_{9,9}$	$\mathbf{B}_{9,10}\mathbf{U}_{10}$	0	0
	0	0	0	0	0	0	0	0	$\mathbf{B}_{10,9}\mathbf{U}_{9}$	$\mathbf{A}_{10,10}$	$\mathbf{B}_{10,11}\mathbf{U}_{11}$	$\mathbf{B}_{10,12}\mathbf{U}_{12}$
_	0	0	0	0	0	0	$\mathbf{B}_{11,7}\mathbf{U}_7$	0	0	$\mathbf{B}_{11,10}\mathbf{U}_{10}$	${f A}_{11,11}$	$\mathbf{B}_{11,12}\mathbf{U}_{12}$
	0	0	0	0	0	0	0	0	0	$\mathbf{B}_{12,10}\mathbf{U}_{10}$	${f B}_{12,11}{f U}_{11}$	$\mathbf{A}_{12,12}$
_												

Components corresponding to the 1-st level cooperation



3	4	5	6	7	8	9	10	11	12

	1	2	3	4	5	6	7	8	9	10	11	12
1	*	1										
2	1	*	1		1							
3		1	*	1								
4			1	*	1	1						
5		1		1	*	1		1				
6				1	1	*	1					
7						1	*	1	1		1	
8					1		1	*	1			
9							1	1	*	1		
10									1	*	1	1
11							1			1	*	1
12										1	1	*

Parity part of the generator matrix



- Components corresponding to the 1-st level cooperation
- Components corresponding to higher level cooperation
 - Can be divided into different groups: each group represents a cycle



Parity part of the generator matrix



Components corresponding to the 1-st level cooperation

- Components corresponding to higher level cooperation
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 - Nodes on each vertical edge have the same cooperation level
 - Nodes on each horizontal edge has the same component matrix that transfer a message into a vector of cross parities



Cooperation Graphs

- The aforementioned matrix is referred to as a *cooperation matrix*
 - Construct the cooperation graph from the cooperation matrix
 - Each cycle represents a pair of edges/triangles, with an arrow pointing from one to the other, labelled with the cooperation level



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Cooperation graph depicts the cooperation of information, i.e., how the information is coupled in the coded DSN



Compatible graphs

>

Cooperation graphs that adopts the following construction

• Higher level cooperation: $\mathbf{B}_{i,j}\mathbf{V}_{j;l}$ ($k_i \times \eta_{j;l}, \eta_{j;l} \times r_j$)

•
$$\mathbf{E}_{i;l;t}$$
 $(k_i \times \gamma_{i;t})$, $\mathbf{B}_{i,j} = \left[\mathbf{E}_{i;l;t}, \mathbf{0}_{k_i \times (\eta_{j;l} - \gamma_{i;t})}\right]$, $t \in T_{i;l}, j \in Y_{t;i}$

• $T_{i;l}$: indices of all cycles that provide extra parities in the l-th level cooperation of node v_i

$$\mathbf{T}_{i} = \begin{bmatrix} \mathbf{A}_{i,i} & \mathbf{B}_{i} & \mathbf{E}_{i;2} & \dots & \mathbf{E}_{i;L_{i}} \\ \hline \mathbf{U}_{i} \\ \hline \mathbf{V}_{i;2} \\ \vdots \\ \hline \mathbf{V}_{i;L_{i}} \end{bmatrix} \xrightarrow{\mathbf{Z}_{i}} \mathbf{E}_{i;l} = \begin{bmatrix} \mathbf{E}_{i;l;t_{1}} & | \dots & | \mathbf{E}_{i;l;t_{|T_{i;l}|}} \end{bmatrix} \\ \mathbf{B}_{i} = \begin{bmatrix} \mathbf{B}_{i,j_{1}} & | \dots & | \mathbf{B}_{i,j_{|\mathcal{M}_{i}|}} \end{bmatrix} \\ \mathbf{ECC Hierarchy} \xrightarrow{\mathbf{m}_{i} \xrightarrow{\mathbf{B}_{i,j}} \mathbf{m}_{i} \mathbf{B}_{i,j}} \xrightarrow{\mathbf{U}_{j}(\mathbf{V}_{j;l})} \underbrace{\mathbf{m}_{i} \mathbf{B}_{i,j} \mathbf{U}_{j}(\mathbf{m}_{i} \mathbf{B}_{i,j} \mathbf{V}_{j;l})}_{Encoded cross parities} \\ d_{i,0} = r_{i} - \delta_{i} - \sum_{l=2}^{L_{i}} \eta_{i;l}, d_{i,1} = r_{i} + \sum_{v_{j} \in \mathcal{M}_{i}} \delta_{j}, \\ d_{i,2} = r_{i} + \sum_{v_{j} \in \mathcal{M}_{i}} \delta_{j} + \sum_{2 \leq l' \leq l, t \in T_{i;l'}} \gamma_{i;t} \end{bmatrix}$$

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Outline

- Introduction
 - Motivation and Model
 - Existing work
- Cooperative Data Protection
 - ECC hierarchy
 - Single-level cooperation
- Multi-level cooperation
 - Cooperation graphs and compatible graphs
 - Construction over compatible graphs

Conclusion

Conclusion

Main contributions

- A hierarchical coding framework that provides topology-aware cooperative data protection for DSNs
- The framework supports scalability and flexibility
- Our scheme achieves faster recovery speed and corrects more flexible erasure patterns

Follow-up work (already done)

- > Analysis of recoverable erasure patterns
- Algorithms that search for cooperation graphs in a network

Future work

- Codes that support non-locally decodable neighboring nodes
- Error correction for latency-sensitive devices at the edge

Thank you!

Q&A

^{*} For further reference

^[1] S. Yang et al., "Topology-aware cooperative data protection in blockchain-based decentralized storage networks", in Proc. IEEE ISIT, Los Angeles, CA, USA, Jun. 2020, pp. 622-627

^[2] S. Yang et al., "Hierarchical coding for cloud storage: topology-adaptivity, scalability, and flexibility", submitted to IEEE TIT, 2020, available on https://arxiv.org/abs/2009.09146